## The Rabin-Miller Primality Test

## Fermat Pseudoprimes; The Fermat Primality Test

Fermat's Little Theorem allows us to prove that a number is composite without actually factoring it.

## Fermat's Little Theorem (alternate statement): If $a^{n-1} \not \equiv 1(\bmod n)$

 for some $a$ with $a \neq 0(\bmod n)$, then $n$ is composite.This statement is absolute: There are no exceptions.

Unfortunately, the inverse statement is not always true.

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Inverse to Fermat's Little Theorem (not always true): If a}\mp@subsup{a}{}{n-1}\equiv
(mod}n)\mathrm{ for some }a\mathrm{ with }a\not=0(\operatorname{mod}n)\mathrm{ , then }n\mathrm{ is prime.
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Some counterexamples:

$$
\begin{aligned}
& 2^{340} \equiv 1(\bmod 341), \text { but } 341=11 \cdot 31 \text { is composite, and } \\
& 5^{560} \equiv 1(\bmod 561), \text { but } 561=3 \cdot 11 \cdot 17 \text { is composite }
\end{aligned}
$$

We say that 341 is a Fermat pseudoprime (to the base 2), and 561 is a Fermat pseudoprime to the base 5.

It is even possible for $a^{n-1} \equiv 1(\bmod n)$ to hold for every $a$ with $\operatorname{gcd}(a, n)=1$, and still have $n$ be composite.

This occurs if $n$ is a Carmichael number (also called an absolute Fermat pseudoprime). A Carmichael number is a Fermat pseudoprime to any base $a$ with $\operatorname{gcd}(a, n)=1$.

Carmichael numbers are fairly rare: There are only seven less than 10000: 561, 1105, 1729, 2465, 2821, 6601, 8911

In fact, there are only 585,355 Carmichael numbers less than $10^{17}$.
Given a randomly chosen odd integer $n$ less than $10^{17}$, the probability that $n$ is a Carmichael number is only a little over $10^{-11}$ (about one in one hundred billion).

For a randomly chosen odd integer $n$ with 100 to 300 digits, the probability that $n$ is a Carmichael number appears to be exceedingly low (for practical purpose, zero).

If $n$ is composite and not a Carmichael number, then there are at most $\varphi(n) / 2$ values of $a(1 \leq a<n)$ for which $a^{n-1} \equiv 1(\bmod n)$.

Let $n$ be any odd integer, other than a Carmichael number.
Say we choose 50 random integers $a$ and compute that each satisfies $a^{n-1} \equiv 1(\bmod n)$.
The probability that this would occur if $n$ is composite is at most $2^{-50} \approx 10^{-15}$.
So we can say with reasonable certainty that $n$ is prime.

If $n$ is composite and not a Carmichael number, then it is actually possible to have $\varphi(n) / \overline{/ 2}$ values for which $a^{n-1} \equiv 1(\bmod n)$.

For example, take $n=91=7 \cdot 13 . \quad \varphi(n)=6 \cdot 12=72$.
There are 36 values of $a$ with $a^{72} \equiv 1(\bmod 91)$, namely $a=1,3$, $4,9,10,12,16,17,22,23,25,27,29,30,36,38,40,43,48,51$, $53,55,61,62,64,66,68,69,74,75,79,81,82,87,88,90$.

But this is unusual.

For nearly all odd composite integers $n$ (other than Carmichael numbers), $a^{n-1} \equiv 1(\bmod n)$ for far fewer than $\varphi(n) / 2$ values of $a$.

For example, let us look at odd composite integers starting with 10001.

| $\boldsymbol{n}$ | $\boldsymbol{\varphi}(\mathbf{n})$ | No of $\boldsymbol{a}$ with <br> $\left.\boldsymbol{a}^{\boldsymbol{n - 1}} \equiv \mathbf{1 ( m o d} \boldsymbol{n}\right)$ |
| :---: | :---: | ---: |
| 10001 | 9792 | 64 |
| 10003 | 8568 | 36 |
| 10005 | 4928 | 64 |
| 10011 | 6440 | 280 |
| 10013 | 8640 | 16 |
| 10015 | 8008 | 4 |
| 10017 | 5616 | 16 |
| 10019 | 9744 | 4 |
| 10021 | 9100 | 100 |
| 10023 | 6144 | 8 |
| 10025 | 8000 | 32 |
| 10027 | 9720 | 162 |
| 10029 | 6684 | 4 |
| 10031 | 8592 | 4 |
| 10033 | 9828 | 36 |
| 10035 | 5328 | 8 |
| 10041 | 6192 | 4 |
| 10043 | 9020 | 4 |

This means that far fewer than the 50 random values of $a$, mentioned earlier, are typically sufficient to show that an odd integer (not a Carmichael number) is prime, with near certainty.

For a randomly chosen odd integer $n$ with 100 to 300 digits, it appears that if $a^{n-1} \equiv 1(\bmod \mathrm{n})$ for even a single randomly chosen $a$, then $n$ is prime with probability very close to 1 .

Fermat Test for Primality: To test whether n is prime or composite, choose $a$ at random and compute $a^{n-1}(\bmod n)$.
i) If $a^{n-1} \equiv 1(\bmod n)$, declare $n$ a probable prime, and optionally repeat the test a few more times.
ii) If $a^{n-1} \not \equiv 1(\bmod n)$, declare $n$ composite, and stop.

We have seen that the Fermat test is really quite good for large numbers.

One limitation: If someone is supposed to provide us with a prime number, and sends a Carmichael number instead, we cannot detect the deception with the Fermat test.
In any case, we can improve upon the Fermat test at almost no cost.

## Euler Pseudoprimes; The Euler Test

If $n$ is an odd prime, we know that an integer can have at most two square roots, $\bmod n$. In particular, the only square roots of $1(\bmod n)$ are $\pm 1$.

If $a \not \equiv 0(\bmod n), a^{(n-1) / 2}$ is a square root of $a^{(n-1)} \equiv 1(\bmod n)$, so $a^{(n-1) / 2} \equiv \pm 1(\bmod n)$.

> If $a^{(n-1) / 2} \not \equiv \pm 1(\bmod n)$ for some $a$ with $a \not \equiv 0(\bmod n)$, then $n$ is composite.

Euler Test: For a randomly chosen $a$ with $a \neq 0(\bmod n)$, compute $a^{(n-1) / 2}(\bmod n)$.
i) If $a^{(n-1) / 2} \equiv \pm 1(\bmod n)$, declare $n$ a probable prime, and optionally repeat the test a few more times.

If $n$ is large and chosen at random, the probability that $n$ is prime is very close to 1 .
ii) If $a^{(n-1) / 2} \not \equiv \pm 1(\bmod n)$, declare $n$ composite.

This is always correct.

The Euler test is more powerful than the Fermat test.
If the Fermat test finds that $n$ is composite, so does the Euler test.
But the Euler test may find $n$ composite even when the Fermat test fails. Why?

If $n$ is an odd composite integer (other than a prime power), 1 has at least 4 square roots $\bmod n$.
So we can have $a^{(n-1) / 2} \equiv \beta(\bmod n)$, where $\beta \neq \pm 1$ is a square root of 1 . Then $a^{n-1} \equiv 1(\bmod n)$. In this situation, the Fermat Test (incorrectly) declares $n$ a probable prime, but the Euler test (correctly) declares $n$ composite.

We noted earlier that
$2^{340} \equiv 1(\bmod 341)$, even though 340 is composite, and
$5^{560} \equiv 1(\bmod 561)$, even though 561 is composite.
We can compute that
$2^{170} \equiv 1(\bmod 341)$, even though 340 is composite, but $5^{280} \equiv 67 \not \equiv \pm 1(\bmod 561)$, showing that 561 is composite.

We call 341 an Euler pseudoprime to the base 2.
But note that 561 is not an Euler pseudoprime base 5, even though it is a Fermat pseudoprime base 5.

On the whole, there are only about half as many Euler pseudoprimes as Fermat pseudoprimes.

Consider the seven Carmichael numbers less than 10000 .
The Euler test can show that 5 of the 7 numbers are composite.

| $\boldsymbol{n}$ | $\boldsymbol{\varphi}(\boldsymbol{n})$ | No of $\boldsymbol{a}$ with <br> $\boldsymbol{a}^{\boldsymbol{n - 1}} \equiv \mathbf{1}(\mathbf{m o d} \boldsymbol{n})$ | No of $\boldsymbol{a}$ with <br> $a^{(\boldsymbol{n}-\mathbf{1}) / \mathbf{2}} \equiv \pm \mathbf{1}(\bmod \boldsymbol{n})$ |
| :---: | :---: | :---: | :---: |
| 561 | 320 | 320 | 160 |
| 1105 | 768 | 768 | 384 |
| 1729 | 1296 | 1296 | 1296 |
| 2465 | 1792 | 1792 | 1792 |
| 2881 | 2160 | 2160 | 1080 |
| 6601 | 5280 | 5280 | 2640 |
| 8911 | 7128 | 7128 | 1782 |

The integers 1729 and 2465 are called absolute Euler pseudoprimes (by analogy with the absolute Fermat pseudoprimes, i.e., Carmichael numbers).

These are composite odd integers such that $a^{(n-1) / 2} \equiv \pm 1(\bmod n)$ for every $a$ with $\operatorname{gcd}(a, n)=1$.

These number cannot be proven composite with the Euler test (unless we happen to choose an $a$ with $\operatorname{gcd}(a, n)>1$, which is exceedingly unlikely if $n$ is a large integer lacking small prime factors.

There are fewer absolute Euler pseudoprimes than there are Carmichael numbers, but unfortunately absolute Euler pseudoprimes do exist.

## The Rabin-Miller Primality Test

The Euler test improves upon the Fermat test by taking advantage of the fact, if 1 has a square root other than $\pm 1(\bmod n)$, then $n$ must be composite.

If $a^{(n-1) / 2} \equiv \pm 1(\bmod n)$, where $\operatorname{gcd}(a, n)=1$, then $n$ must be composite for one of two reasons:
i) If $a^{n-1} \not \equiv 1(\bmod n)$, then $n$ must be composite by Fermat's Little Theorem
ii) If $a^{n-1} \equiv 1(\bmod n)$, then $n$ must be composite because $a^{(n-1) / 2}$ is a square root of $1(\bmod n)$ different from $\pm 1$.

The limitation of the Euler test is that is does not go to any special effort to find square roots of 1 , different from $\pm 1$. The Rabin-Miller test does do this.

For example, recall the Euler Test declares 341 a probable prime because $2^{170} \equiv 1(\bmod 341)$.

But if we compute $2^{85}(\bmod 341)$, we find $2^{85} \equiv 32(\bmod 341)$. Thus 32 is a square root of $2^{2 \cdot 85} \equiv 2^{170} \equiv 1(\bmod 341)$, different from $\pm 1$, so we would find that 341 is composite.

In the Rabin-Miller test, we write $n-1=2^{s} \cdot m$, with $m$ odd and $s \geq 1$.
We then start by compute $a^{m}(\bmod n)$ using fast exponentiation.
If $a^{m} \equiv \pm 1(\bmod n)$, we declare n a probable prime, and stop.
Why? We know that $a^{n-1} \equiv\left(a^{m}\right)^{2^{s}} \equiv 1(\bmod n)$, and we will not
find a square root of 1 , other than $\pm 1$, in repeated squaring of $a^{m}$ to get $a^{n-1}$.

Otherwise, unless $s=1$, we square $a^{m}(\bmod n)$ to obtain $a^{2 m}$.
If $a^{2 m} \equiv 1(\bmod n)$, we declare $n$ composite, and stop.
Why? $a^{m}$ is a square root of $a^{2 m} \equiv 1(\bmod n)$, different from $\pm 1$.

If $a^{2 m} \equiv-1(\bmod n)$, we declare $n$ a probable prime, and stop.
Why? Just as above, we know that $a^{n-1} \equiv 1(\bmod n)$, and we will not find a square root of 1 , other than $\pm 1$.

Otherwise, unless $s=2$, we square $a^{2 m}(\bmod n)$ to obtain $a^{2^{2} m}$.

If $a^{2^{2} m} \equiv 1(\bmod n)$, we declare $n$ composite, and stop.
Why? We have found a square root of $1(\bmod n)$, different from $\pm 1$, just as above

If $a^{2 m} \equiv-1(\bmod n)$, we declare $n$ a probable prime, and stop.
Why? Just above, we know that $a^{n-1} \equiv 1(\bmod n)$, and we will not find a square root of 1 , other than $\pm 1$.

Otherwise we continue in this manner until either (a) we stop the test, or (b) we have computed $a^{2^{s-1} m}$, and stopped if $a^{2^{s-1} m} \equiv a^{(n-1) / 2} \equiv \pm 1(\bmod n)$.

If we haven't stopped by this point, we declare $n$ composite and stop.

Why? Exactly as with the Euler test.

Let us carry out the Rabin-Miller test on the absolute Euler pseudoprime 1729, using $a=671$.

$$
\begin{array}{cc}
1729-1=1728=2^{6} \cdot 27 . & \text { So } s=6, m=27 . \\
671^{27} \equiv 1084 & (\bmod 1729) \\
671^{27 \cdot 2} \equiv 1084^{2} \equiv 1065 & (\bmod 1729) \\
671^{27 \cdot 2^{2}} \equiv 1065^{2} \equiv 1 & (\bmod 1729)
\end{array}
$$

The test declares $n$ composite, and terminates.

Next we test a much larger integer, $n=972133929835994161$ (also a Carmichael number), using $a=2$.

$$
\begin{array}{rlr}
n-1=2^{4} \cdot 60758370614749635 . \\
2^{60758370614749635} & \equiv 338214802923303483 & (\bmod n) \\
2^{2 \cdot 60758370614749635} & \equiv 338214802923303483^{2} & (\bmod n) \\
& \equiv 332176174063516118 & (\bmod n) \\
2^{2^{2} \cdot 60758370614749635} & \equiv 332176174063516118^{2} & (\bmod n) \\
& \equiv 779803551049098051 & (\bmod n) \\
2^{2^{3} \cdot 60758370614749635} & \equiv 779803551049098051^{2} & (\bmod n) \\
& \equiv 1 & (\bmod n)
\end{array}
$$

The test declares $n$ composite, and terminates.

Next we test an integer that is composite, but not a Carmichael number, $n=2857191047211793$, using $a=1003$.

$$
\begin{array}{rlr}
n-1=2^{4} \cdot 178574440450737 . \\
1003^{178574440450737} & & \equiv 1135781085623492 \\
1003^{2 \cdot 178574440450737} & (\bmod n) \\
& \equiv 1135781085623492^{2} & (\bmod n) \\
& \equiv 84313648747407 & (\bmod n) \\
1003^{2^{2} \cdot 178574440450737} & \equiv 84313648747407^{2} & (\bmod n) \\
& 2321094267189023 & (\bmod n) \\
1003^{2^{3} \cdot 178574440450737} & \equiv 2321094267189023^{2} & (\bmod n) \\
& \equiv 978857874792606 & (\bmod n)
\end{array}
$$

The test declares $n$ composite, and terminates.

Finally we test an integer that is in fact prime, $n=104513$, using $a=$ 3.

$$
\begin{array}{rlr}
n-1=2^{6} \cdot 1633 . \\
& & \\
3^{1633} & \equiv 88958 & \\
3^{2 \cdot 1633} & \equiv 88958^{2} & \equiv 10430 \\
(\bmod n) \\
3^{2^{2} \cdot 1633} & \equiv 10430^{2} & \equiv 91380 \\
(\bmod n) \\
3^{2^{3} \cdot 1633} & \equiv 91380^{2} & \equiv 29239 \\
(\bmod n) \\
3^{2^{4} \cdot 1633} & \equiv 29239^{2} & \equiv 2781 \\
3^{2^{5} \cdot 1633} & \equiv 2781^{2} & \equiv-1
\end{array} \quad(\bmod n)
$$

The test concludes that $n$ is a probable prime. We might perform a few more tests before we are convinced that $n$ is in fact prime.

Like the Fermat and Euler tests, the Rabin-Miller test has psuedoprimes (choices of $a$ for which the test declares a composite integer to be a probable prime).

Rabin-Miller pseudoprimes are called strong pseudoprimes.
There are fewer strong pseudoprimes than Fermat or Euler pseudoprimes.

More importantly, there are no Rabin-Miller absolute pseudoprimes (as we had absolute Fermat and Euler absolute pseudoprimes).

For any odd composite integer $n$, there are at most $\varphi(n) / 4$ integers $a$ $(1 \leq a<n, \operatorname{gcd}(a, n)=1)$ for which the Rabin-Miller test declares $n$ prime.

In practice, the number of strong pseudoprimes is usually far, far less than $\varphi(n) / 4$, if $n$ is large.

There are a number of other primality tests, but the Rabin-Miller test is the one most commonly used.

