

Internal Languages for Higher Categories

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Outline

- ① Internal languages for categories
- ② Internal languages for higher categories
- ③ Research problems

The multiverse of mathematics

Theorem (“Gödel’s incompleteness theorem”)

No strong, sensible formal system has a unique model.

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Possible reactions:

- ① **Oh no!** I thought I was studying one thing, but all this time I might have been studying something else entirely.
- ② **Hey cool!** All of my theorems are **more general** than I thought they were! That seems likely to be really useful.

Nonstandard models

Example

- Formal system: Grade-school arithmetic
- Classical model: The real numbers
- Nonstandard models: Other fields

Nonstandard models

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- **Classical model:** The real numbers
- **Nonstandard models:** Other fields

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- **Classical model:** Euclidean geometry
- **Nonstandard models:** Non-Euclidean geometries

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Example

- Formal system: Zermelo–Fraenkel set theory
- Classical model: The “real” universe of sets
- Nonstandard models: **Alternative worlds of mathematics**

Type theory

Classical set-theoretic foundations are not ideal for working with alternative models. A better choice is **type theory**.

What you need to know about type theory for this talk

- Its basic objects are **types**, which are kind of like sets.
- Types contain **terms**, which are kind of like elements of sets.
- “ $a : A$ ” means “the term a belongs to the type A ”.
- “ $(x : A) \vdash (b : B)$ ” means “assuming that x is a variable belonging to the type A , the term b belongs to the type B ”.

Categorical semantics

The “more general contexts” where we can interpret type theory are **categories** with appropriate structure.

Type theory	Model in a category \mathbf{C}
Type $(A : \text{Type})$	Object $\llbracket A \rrbracket \in \mathbf{C}$
Term $(x : A) \vdash (b : B)$	Morphism $\llbracket A \rrbracket \xrightarrow{\llbracket b \rrbracket} \llbracket B \rrbracket$

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Slogan

Objects of any category “look internally” like sets.

Categorical semantics: example

Example

In mathematics built on type theory, a **group** is a type G with terms

$$\begin{aligned} & \vdash (e : G) \\ (x : G), (y : G) & \vdash (x \cdot y : G) \\ (x : G) & \vdash (x^{-1} : G) \end{aligned}$$

satisfying the appropriate axioms.

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Interpreted internally, we get

- In sets: a group.
- In topological spaces: a topological group.
- In manifolds: a Lie group.
- In sheaves: a sheaf of groups.
- In rings^{op}: a Hopf algebra.

Categorical semantics: really useful!

Theorem (proven in type theory)

The inversion map of a group is unique.

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The antipode of a Hopf algebra is unique.

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The antipode of a Hopf algebra is unique.

NB: Not all classical reasoning is valid internally, e.g. the **law of excluded middle** and the **axiom of choice** usually fail. But there's still a lot left.

Nonclassical axioms

- So far, we've been talking about interpreting **classically true** statements in more general contexts.
- We can also consider **classically false** statements (that is, statements which are false in the “originally intended” model) which are true in other interesting contexts.

Nonclassical axioms

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Examples

- All functions $\mathbb{R} \rightarrow \mathbb{R}$ are continuous.
- All functions $\mathbb{N} \rightarrow \mathbb{N}$ are computable.
- There exist real “infinitesimals” with which we can do calculus (*synthetic differential geometry*).

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Higher categories

A **higher category** is a category together with “higher morphisms” or “higher homotopies” between its morphisms. Thus it has:

- Objects A, B, \dots
- Morphisms $A \rightarrow B$,
- 2-morphisms $A \begin{array}{c} \curvearrowright \\ \downarrow \\ \curvearrowleft \end{array} B$
- 3-morphisms $A \begin{array}{c} \curvearrowright \\ \left(\leftarrow \right) \\ \curvearrowleft \end{array} B$
- \dots

We will consider only **$(\infty, 1)$ -categories**, where all the higher morphisms are “invertible”.

Higher categories

Examples

- Topological spaces
 - Simplicial sets
 - Chain complexes
 - Categories
 - Spectra (stable spaces)
 - Simplicial sheaves
- } the “original model”, analogous to sets

Homotopy type theory

The **absolutely magical fact** is that one of the most natural type theories is perfectly adapted to models in $(\infty, 1)$ -categories.

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$(x : A) \vdash (p : \text{Id}_B(b, b'))$	2-morphism $\llbracket A \rrbracket \begin{array}{c} \xrightarrow{\quad} \\ \downarrow \\ \xrightarrow{\quad} \end{array} \llbracket B \rrbracket$
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Slogan

Objects of an $(\infty, 1)$ -category “look internally” like **spaces**, a.k.a. **homotopy types**, a.k.a. **∞ -groupoids**.

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We need to make the correspondence between type theory and higher categories precise (there are coherence issues).

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What we have:

- An almost-complete model in any $(\infty, 1)$ -topos (Lumsdaine–Warren and others).
- A complete model in the “standard” case of simplicial sets (Voevodsky).
- Complete models in a small class of presheaf $(\infty, 1)$ -toposes (Shulman).

Problem 2

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Which parts of classical homotopy theory are still true internally?

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Some things we know:

- The law of excluded middle, the axiom of choice, all fail in general (just as for 1-categories).
- “Whitehead’s theorem” fails: a homotopy type is not determined by its homotopy groups (Lurie).
- But $\pi_1(S^1) = \mathbb{Z}$ (Shulman).

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Open Problem

Are all the homotopy groups of spheres determined by type theory?
Or could they be different in different categories?

Problem 3

Question

What interesting nonclassical axioms can hold in higher categories?

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Example

A **cohesive $(\infty, 1)$ -topos** is one whose objects are homotopy types equipped with some sort of topological or smooth structure. Its internal type-theory admits **higher modalities**, using which one can describe large chunks of differential cohomology and gauge theory (Schreiber–Shulman).

Thanks!